

prolog :-in(2016).

# Prolog

- Based on first-order predicate calculus
  - Allow Horn clauses (a subset of predicate calculus) only
  - Execution of a Prolog program is an application of theorem proving
- Original motivation: study of mechanical theorem proving
- Developed in 1970 by Colmerauer & Roussel (Marseilles) and Kowalski (Edinburgh) + others.
  - The name Prolog is taken from PROgramming in LOGic.
- Used in artificial intelligence, databases, expert systems.
  - Non numeric applications
- Popular in Europe
- Adopted by Japan for its 5th Generation Computer program in 80's
- Viewed as competitor for Lisp

# Prolog is a declarative programming language

- Express programs in a form of symbolic logic and use a logical inferencing process to produce results
- Logic programs are declarative rather than procedural
  - only the specification of the desired results are stated rather detailed procedures for producing them

# Run Prolog

- Popular Prolog implementations
  - GNU Prolog, SWI Prolog, ....
- SWI Prolog is installed on Luna machine
  - The command to run swi prolog: pl
  - The command to run gprolog: gprolog
- Edit the program
  - suppose the file sample.pl consists of the following Prolog program:  
`likes(john, mary).`

- Run the program (red part are user inputs)

```
luna:~/440/prolog>gprolog
GNU Prolog 1.2.16
By Daniel Diaz
Copyright (C) 1999-2002 Daniel Diaz
| ?- consult('sample.pl').
compiling /fac2/jlu/440/prolog/sample.pl for
byte code...
/fac2/jlu/440/prolog/sample.pl compiled, 47
lines read - 5711 bytes written, 33 ms
(10 ms) yes
| ?- likes(john, mary).
yes

| ?- likes(john, Whom).
Whom = mary
yes

| ?- likes(john, george).
No
```

# Prolog program and query

- In Prolog you define a set of facts and rules, then you can query based on your definitions;
- The basic unit of your program is predicate;
- A predicate consists of a name and a sequence of arguments
  - `likes(john, mary).`
  - “likes” is the predicate name, `john` and `mary` are two arguments of the predicate.
- Once you enter this fact into the system, you can ask two types of questions:
  - true/false question: `?-like(john, mary) .` does john like mary?
  - what question: `?-likes(john, X) .` whom does john like? The X here is a variable.

## Components of Prolog programs--facts

- A fact has the form  $p(t_1, \dots, t_n)$ .

- Examples:

```
likes(john, mary) .
```

```
parent(adam, bill) .
```

- A fact is part of a program.

- Syntactic conventions

- It is case sensitive:

- Predicate and atom names must begin with a lowercase letter;

- Every statement end with a dot “.”.

## Facts with variables

- Variables starts with uppercase letters.
- Facts can have variables

```
likes(john, mary).
```

```
likes(dwight, X).
```

```
    % dwight likes anybody (or anything).
```

```
    %  $\forall X$ . likes(dwight, X).
```

```
| ?- likes(dwight, mary).
```

```
yes
```

```
| ?- likes(dwight, john).
```

```
yes
```

```
| ?- likes(dwight, a).
```

```
yes
```

# Rules

- The second type of statements is the rule.

likes(john, X):-likes(mary, X).

- The symbol “:-” means “if”.

- This rule means that if mary likes X, then john also likes X.

- Rules make use of variables.  $\forall X. (\text{likes}(\text{mary}, X) \rightarrow \text{likes}(\text{john}, X))$ .

- Running example

likes(john, mary). likes(dwight, X). likes(mary, sue).

likes(john, X):-likes(mary, X).

- | ?- likes(john, X).

- X = mary ? ;

- X = sue

- yes

- | ?-

- Prolog will give you one answer for X;

- If you hit return it will stop looking for more answers;

- If you type “;”, it will look for another answer;

- Prolog marks the database where it finds answers in order to start looking again.



## Rule: clause body can have multiple predicates

father(X,Y) :- parent(X,Y), male(X).

- The antecedent and consequent are in reverse order to that normally found in logic:

parent(X,Y)  $\wedge$  male(X)  $\rightarrow$  father(X,Y)

- the consequent is written first and called the head of the rule;
- the antecedent is called the body;
- Conjunction (and) is written as ", ".


parent(bill, dwight).

male(bill).

father(X,Y):-parent(X,Y), male(X).

| ?- father(bill, X).

X = dwight

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## General form of a rule (Horn Clause)

- $A :- B_1, B_2, \dots, B_n.$   
meaning A is true if B1 and B2 and ... Bn are all true
- Clause: a disjunction of literals
- Horn clause: if it contains at most one positive literal.

$$\begin{aligned} A & :- \neg B_1, B_2, \dots, B_n \\ & \equiv B_1 \wedge B_2 \wedge \dots \wedge B_n \rightarrow A \\ & \equiv \neg(B_1 \wedge B_2 \wedge \dots \wedge B_n) \vee A \\ & \equiv \neg B_1 \vee \neg B_2 \vee \dots \vee \neg B_n \vee A \end{aligned}$$

## Evaluation--unification

- When the interpreter is given a query, it tries to find facts that match the query.

```
likes(john, mary).  
likes(john, sue).  
| ?- likes(john, X).
```

- Iterates in order through the program's "likes" clauses
  - First match is `likes(john, mary)`, so it returns `X=mary`.
  - If you hit semicolon ";", it continues and find the next match. So it returns `X=sue`.

- Unification

```
likes(john, mary).  
|   |  
likes(john, X).
```

- Unification is a two-way match

```
likes(Y, mary).  
|   |  
likes(john, X).
```

```
likes(Y, mary).
```

```
| ?- likes(john, X).
```

```
X=mary
```

## Evaluation—backward chaining

- If no outright facts are available, it attempts to satisfy all rules that have the fact as a conclusion.

- For example:

parent(bill, dwight). male(bill).

father(X,Y):-parent(X,Y), male(X).

| ?- trace.

| ?- father(bill, W).

1 1 Call: father(bill,\_16) ?

2 2 Call: parent(bill,\_16) ?

X=bill

2 2 Exit: parent(bill,dwight) ?

Y=dwight

3 2 Call: male(bill) ?

3 2 Exit: male(bill) ?

1 1 Exit: father(bill,dwight) ?

W = dwight

yes

## Recursive rules

- How to define the ancestor predicate?

```
ancestor(X,Y):-parent(X,Y).
```

```
ancestor(X,Y):-parent(X, U), parent(U, Y).
```

```
ancestor(X,Y):-parent(X, U), parent(U, V),  
parent(V, W), parent(W, Y).
```

... ..

- Prolog is not a pure logic programming language.

- Solution

```
ancestor(X,Y):-parent(X,Y).
```

```
ancestor(X,Z):-parent(X,Y), ancestor(Y,Z).
```

```
parent(bill, dwight).
```

```
parent(dwight, ryan).
```

```
parent(dwight, samantha).
```

```
parent(dwight, sue).
```

```
| ?-ancestor(bill, X).
```

```
X=dwight;
```

```
X=ryan;
```

```
X=samantha;
```

```
X=sue;
```

- Another solution?

```
– ancestor(X,Y):-parent(X,Y).
```

```
– ancestor(X,Z):-ancestor(X, Y), parent(Y,Z).
```



- The correct program:

edge(a,b).

edge(b,c).

edge(c,d).

edge(d,e).

edge(b,e).

edge(d,f).

path(X,Y):-edge(X,Z), path(Z,Y).

path(X,X).

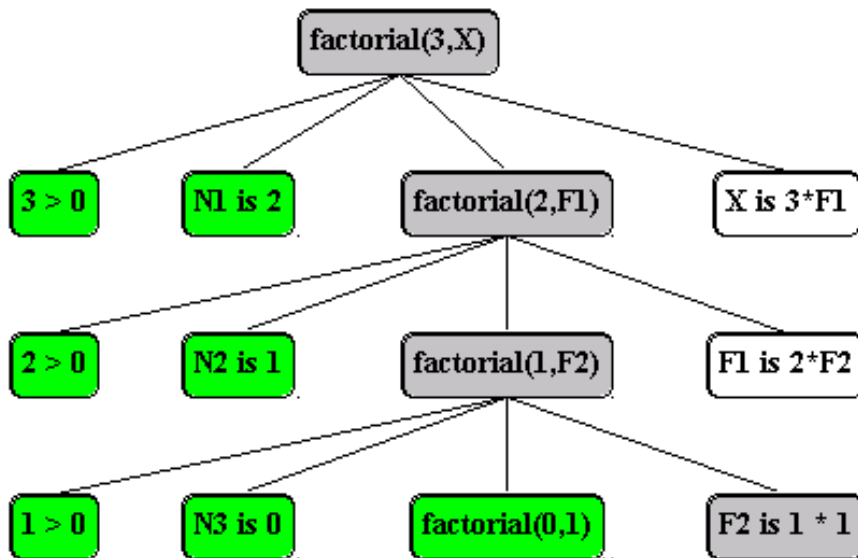
# Evaluation: depth first search

$f(0, 1)$ .

$f(N, F): -N > 0,$

$N1$  is  $N-1, f(N1, F1),$

$F$  is  $N * F1.$



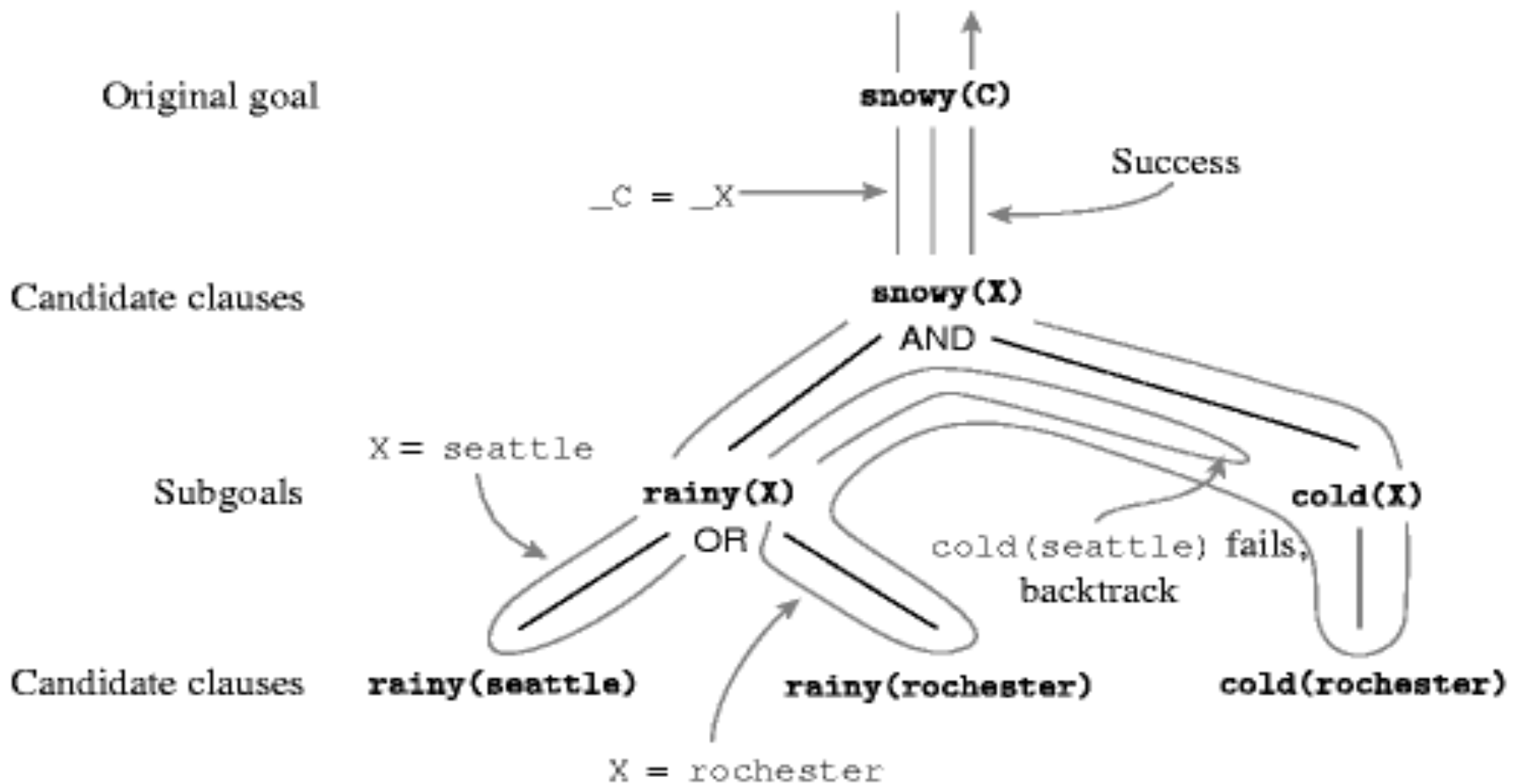
| ?- f(3,X).

- 1 1 Call: f(3,\_16) ?
- 2 2 Call: 3>0 ?
- 2 2 Exit: 3>0 ?
- 3 2 Call: \_113 is 3-1 ?
- 3 2 Exit: 2 is 3-1 ?
- 4 2 Call: f(2,\_138) ?
- 5 3 Call: 2>0 ?
- 5 3 Exit: 2>0 ?
- 6 3 Call: \_190 is 2-1 ?
- 6 3 Exit: 1 is 2-1 ?
- 7 3 Call: f(1,\_215) ?
- 8 4 Call: 1>0 ?
- 8 4 Exit: 1>0 ?
- 9 4 Call: \_267 is 1-1 ?
- 9 4 Exit: 0 is 1-1 ?
- 10 4 Call: f(0,\_292) ?
- 10 4 Exit: f(0,1) ?
- 11 4 Call: \_320 is 1\*1 ?
- 11 4 Exit: 1 is 1\*1 ?
- 7 3 Exit: f(1,1) ?
- 12 3 Call: \_349 is 2\*1 ?
- 12 3 Exit: 2 is 2\*1 ?
- 4 2 Exit: f(2,2) ?
- 13 2 Call: \_16 is 3\*2 ?
- 13 2 Exit: 6 is 3\*2 ?
- 1 1 Exit: f(3,6) ?



# Evaluation: backtracking search

```
rainy(seattle).  
rainy(rochester).  
cold(rochester).  
snowy(X) :- rainy(X), cold(X)
```



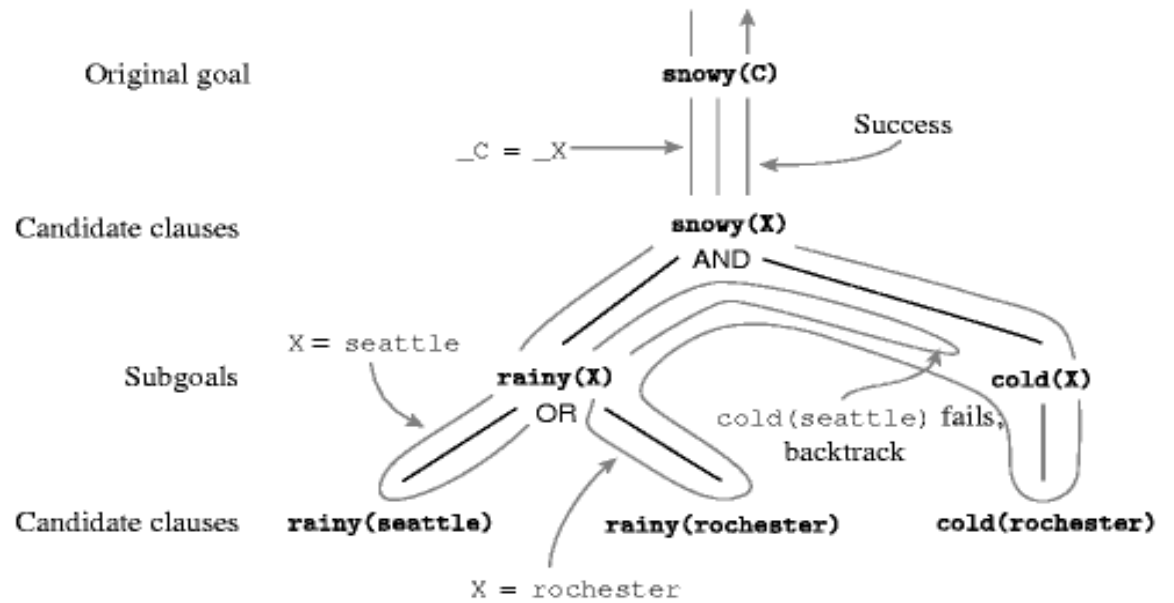
```

rainy(seattle).
rainy(rochester).
cold(rochester).
snowy(X) :- rainy(X), cold(X)

```

| ?- snowy(C).

- 1 1 Call: snowy(\_16) ?
- 2 2 Call: rainy(\_16) ?
- 2 2 Exit: rainy(seattle) ?
- 3 2 Call: cold(seattle) ?
- 3 2 Fail: cold(seattle) ?
- 2 2 Redo: rainy(seattle) ?
- 2 2 Exit: rainy(rochester) ?
- 3 2 Call: cold(rochester) ?
- 3 2 Exit: cold(rochester) ?
- 1 1 Exit: snowy(rochester) ?



C = rochester

## Review

- Syntax: facts, rules. Horn clause

$f(0, 1).$

$f(N, F):-N>0,$

$N1 \text{ is } N-1, f(N1, F1),$

$F \text{ is } N * F1.$

- Evaluation

– Unification

– Backward chaining

– Recursion

– Depth first search

– Backtracking

## Datalog and Prolog

- What we've seen so far is called Datalog: “databases in logic.”
- Prolog is “programming in logic.” It goes a little bit further by allowing complex terms, including records, lists and trees.
- These complex terms are the source of the hard thing about Prolog, “unification.”

`at_uwindsor(student(128327, 'Spammy K', date(2, may, 1986))).`

`at_uwindsor (student(126547, 'Blobby B', date(15, dec, 1985))).`

`at_uwindsor (student(456591, 'Fuzzy W', date(23, aug, 1966))).`

- It is no longer a simple fact `at_uwindsor(snammy)`.

at\_uwindsor(student(128327, 'Spammy K', date(2, may, 1986))).

at\_uwindsor (student(126547, 'Blobby B', date(15, dec, 1985))).

at\_uwindsor (student(456591, 'Fuzzy W', date(23, aug, 1966))).

| ?- at\_uwindsor(student(IDNum, Name, date(Day,Month,Year))), Year < 1983.

Day = 23

IDNum = 456591

Month = aug

Name = 'Fuzzy W'

Year = 1966

at\_uwindsor(student(128327, 'Spammy K',date(2, may, 1986))).

at\_uwindsor (student(126547, 'Blobby B', date(15, dec, 1985))).

at\_uwindsor (student(456591, 'Fuzzy W', date(23, aug, 1966))).

| ?- at\_uwindsor(student(\_, Name, date(\_,\_,Year))), Year < 1983.

Name = 'Fuzzy W'

Year = 1966

| ?- at\_uwindsor(Person), Person=student(\_,\_,Birthday), Birthday=date(\_,\_,Year), Year < 1983.

Birthday = date(23,aug,1966)

Person = student(456591,'Fuzzy W',date(23,aug,1966))

Year = 1966

## List

- data structure built into prolog.
  - `[]` % the empty list
  - `[1]`
  - `[1,2,3]`
- `|` separate the head and tail of a list
  - `[a | b]` adds the first argument `a` to the list `b`.
- `[[1,2], 3]` % a list can be heterogeneous.
- All the following are the same as `[a, b, c]`
  - `[a, b, c|[]]`
  - `[a, b| [c]]`
  - `[a | [b, c]]`

## Decompose a list

- How to get the first element of a list? (car in scheme)

– `first(X, List) :- ...?`

- Solution(s)

```
first(X, List) :- List=[X|Xs].
```

```
first(X, [X|Xs]).
```

```
first(X, [X|_]).
```

– eliminate the single-use variable Xs.

Query: `first(8, [7,8,9]).`

Answer: no

Query: `first(X, [7,8,9]).`

Answer: X=7

How to define rest of a list? `rest(Xs, List):-?`

```
rest(Xs, [_|Xs]).
```



# List processing (append)

```
append([],Ys): return Ys
```

```
append([X|Xs],Ys): return [X | append(Xs,Ys)]
```

- In Prolog there are no return values. Rather, the return value is a third argument: `append(Xs,Ys,Result)`.
- This is a constraint saying that `Result` must be the append of the other lists.
- Any of the three arguments may be known (or partly known) at runtime. We look for satisfying assignments to the others.
- `append(Xs,Ys,Result)` is true if `Xs` and `Ys` are lists and `Result` is their concatenation (another list).
- Try this:

```
append([], Ys, Ys) .
```

```
append([X|Xs], Ys, Result) :- ... ?
```

```
append([], Ys, Ys) .
```

```
append([H|T], L, [H|L2]) :- append(T, L, L2) .
```

```
| ?- append([1,2], [3,4,5], X) .
```

```
X = [1,2,3,4,5]
```

## Run the program backwards

- In Prolog, once you've written it, you can also run it backwards!

```
| ?- append([1,2], Y, [1,2,3,4,5,6]).
```

```
Y = [3,4,5,6]
```

```
| ?- append(A,B,[1,2,3]).
```

```
A = [ ],
```

```
B = [1,2,3] ;
```

```
A = [1],
```

```
B = [2,3] ;
```

```
A = [1,2],
```

```
B = [3] ;
```

```
A = [1,2,3],
```

```
B = [ ] ;
```

```
no
```

# List processing--length

len([],0).

len([H|T], X) :-

len(T, Y),

X is Y + 1.

| ?- len([1,2,3],X).

1 1 Call: len([1,2,3],\_22) ?

2 2 Call: len([2,3],\_91) ?

3 3 Call: len([3],\_115) ?

4 4 Call: len([],\_139) ?

4 4 Exit: len([],0) ?

5 4 Call: \_167 is 0+1 ?

5 4 Exit: 1 is 0+1 ?

3 3 Exit: len([3],1) ?

6 3 Call: \_196 is 1+1 ?

6 3 Exit: 2 is 1+1 ?

2 2 Exit: len([2,3],2) ?

7 2 Call: \_22 is 2+1 ?

7 2 Exit: 3 is 2+1 ?

1 1 Exit: len([1,2,3],3) ?

X = 3

yes

# List processing: member

- `member(X,Y)` should be true if `X` is any object, `Y` is a list, and `X` is a member of the list `Y`.

```
member(X, [X|_]).    % same as first(..).
```

```
member(X, [H|T]) :- member(X, T).
```

```
?- member(3, [1,2,3]).    %member test
```

yes

```
?- member(X, [1,2,3]).    % member generator
```

```
X = 1 ;
```

```
X = 2 ;
```

```
X = 3 ;
```

```
No
```

```
?-member(2, X).    % list generator
```

```
X = [2|_] ? ;
```

```
X = [_ ,2|_] ? ;
```

```
X = [_ ,_ ,2|_] ? ;
```

```
X = [_ ,_ ,_ ,2|_] ?
```

```
...
```

- Query: `member(7, List), member(8, List), length(List, 3)`.
  - Answer: doesn't terminate
- Query: `length(List, 3), member(7, List), member(8, List)`.
  - Answer: `List=[7, 8, X]` ;  
`List=[7, X, 8]` ;  
`List=[8, 7, X]` ;  
`List=[X, 7, 8]` ;  
`List=[8, X, 7]` ;  
`List=[X, 8, 7]`
- Again, only “theoretically” Prolog does not depend on the orders of predicates and clauses.

## Example: test subset relationship

```
subset([ ],Y).
```

```
subset([A|X],Y) :- member(A,Y), subset(X,Y).
```

```
| ?- subset ([a, b, c], [a, c, d, b]).
```

true ?

## Example: nth element

```
find_nth([H|T], 1, H).
```

```
find_nth([H|T], N, X):- N1 is N-1, find_nth(T,N1,X).
```

```
| ?- find_nth([a, b, c, d], 2, X).
```

X = b ?

## Example: what are the changes for a dollar?

% suppose there are half-dollars, quarters, dimes, nickels. P is the purchase	?- change([H,Q,D,N,P]) .	?- change([H, Q,D,N,20]).
change ([H,Q,D,N,P]) :-	D = 0	D = 0
member(H,[0,1,2]), // Half-dollars	H = 0	H = 0
member(Q,[0,1,2,3,4]), // quarters	N = 0	N = 16
member(D,[0,1,2,3,4,5,6,7,8,9,10]) , // dimes	P = 100	Q = 0 ? ;
member(N, [0,1,2,3,4,5,6,7,8,9,10, // nickels	D = 0	D = 1
11,12,13,14,15,16,17,18,19,20]),	H = 0	H = 0
S is 50*H + 25*Q +10*D + 5*N,	N = 1	N = 14
S =< 100, P is 100-S.	P = 95	Q = 0 ? ;
	Q = 0 ? ;	
	D = 0	D = 2
	H = 0	H = 0
	N = 2	N = 12
	P = 90	Q = 0 ?
	Q = 0 ?	
	...	



# Insertion sort

isort([ ],[ ]).

isort([X|UnSorted], AllSorted) :-

isort(UnSorted, Sorted),

insert(X, Sorted, AllSorted).

insert(X, [ ], [X]).

insert(X, [Y|L], [X, Y|L]) :- X =< Y.

insert(X, [Y|L], [Y|IL]) :-

X > Y, insert(X, L, IL).

| ?- isort([1, 3, 2], X).

1 1 Call: isort([1,3,2],\_22) ?

2 2 Call: isort([3,2],\_91) ?

3 3 Call: isort([2],\_115) ?

4 4 Call: isort([],\_139) ?

4 4 Exit: isort([],[]) ?

5 4 Call: insert(2,[],\_165) ?

5 4 Exit: insert(2,[],[2]) ?

3 3 Exit: isort([2],[2]) ?

6 3 Call: insert(3,[2],\_194) ?

7 4 Call: 3=<2 ?

7 4 Fail: 3=<2 ?

7 4 Call: 3>2 ?

7 4 Exit: 3>2 ?

8 4 Call: insert(3,[],\_181) ?

8 4 Exit: insert(3,[],[3]) ?

6 3 Exit: insert(3,[2],[2,3]) ?

2 2 Exit: isort([3,2],[2,3]) ?

9 2 Call: insert(1,[2,3],\_22) ?

10 3 Call: 1=<2 ?

10 3 Exit: 1=<2 ?

9 2 Exit: insert(1,[2,3],[1,2,3]) ?

1 1 Exit: isort([1,3,2],[1,2,3]) ?

## A more declarative sort program

```
sort(List, Sorted):-permutation(List, Sorted), is_sorted(Sorted).
```

```
is_sorted([ ]).
```

```
is_sorted([ _ ]).
```

```
is_sorted([X,Y|T]):-X<=Y, is_sorted([Y|T]).
```

```
perm([ ], [ ]).
```

```
perm(List, [H|Perm]) :-takeout(H, List, Rest),  
    perm(Rest, Perm).
```

```
takeout(X, [X|Xs], Xs) .
```

```
takeout(X, [Y|Xs], [Y|Ys]):-takeout(X,Xs,Ys).
```

# Summary of Prolog Programs

- Program = a bunch of axioms
- Run your program by:
  - Enter a series of facts and declarations
  - Pose a query
  - System tries to prove your query by finding a series of inference steps
- “Philosophically” declarative
- Actual implementations are deterministic
- Many concepts not covered
  - Negation
  - Cut.

Questions?